A Quantitative Approach for Refactoring NFV-based Mobile Core Networks

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- *Motivation*: The network function virtualization (NFV) technology can be used to **improve scalability** and **increase resources utilization**.

- *Problem*: When we apply NFV to mobile core networks (**4G** or **5G**) *in the straightforward*, the **scaling** is **inefficient** (**low resource utilization**) even though the entities can **scale out or in** according to the current load.

- *Main Idea and goal*: To **refactor** mobile core networks. *Phase 1*: To **split** the decomposable entities into smaller network functions to **increase resource utilization**. *Phase 2:* To **merge** selected network functions to **shorten control signaling delay**.

- *Main Contribution:* We **design two quantitative indicators** to evaluate the impacts of **merging** two consecutive functions in the generated strings. The problem to select the network functions to be merged is **reducible to a string matching problem**.

Mobile All-IP Networking LAB

Fooling AI with AI: An Accelerator for Adversarial Attacks on Deep Learning Visual Classification
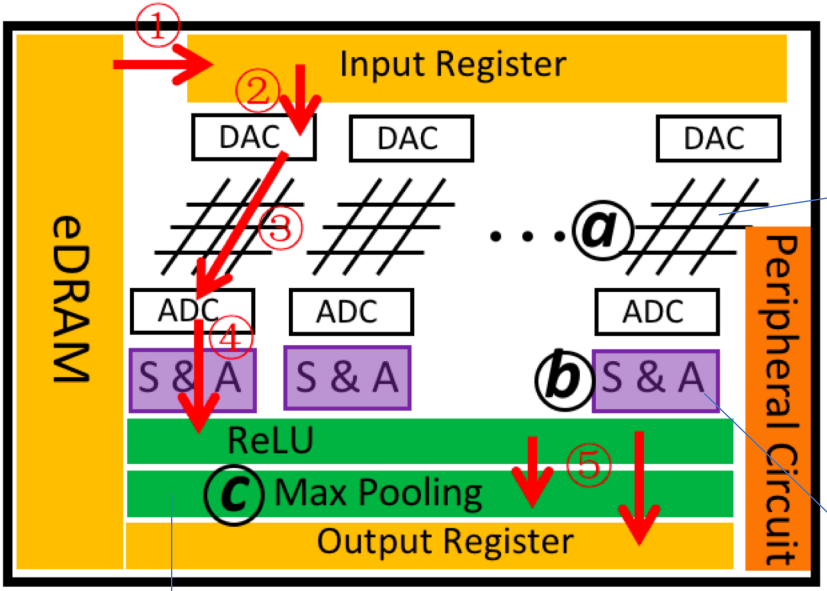
Haoqiang Guo   Lu Peng
Jian Zhang   Fang Qi
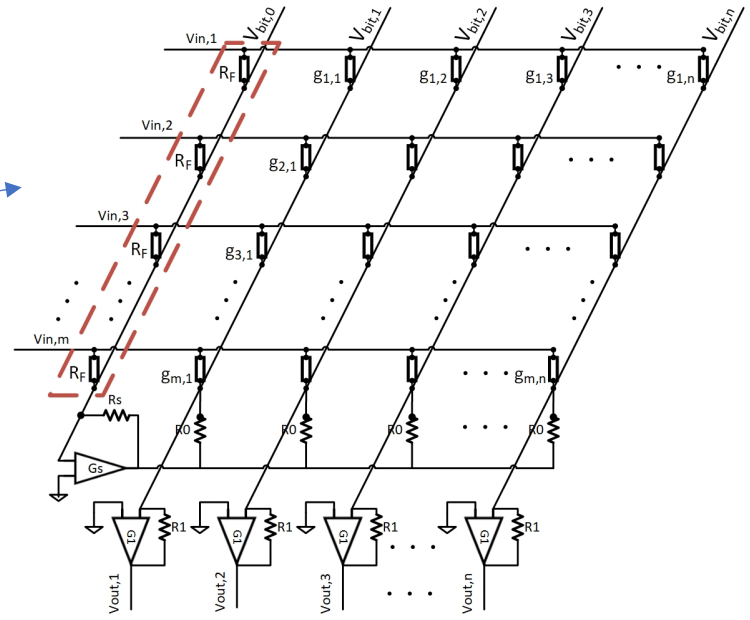
Lide Duan

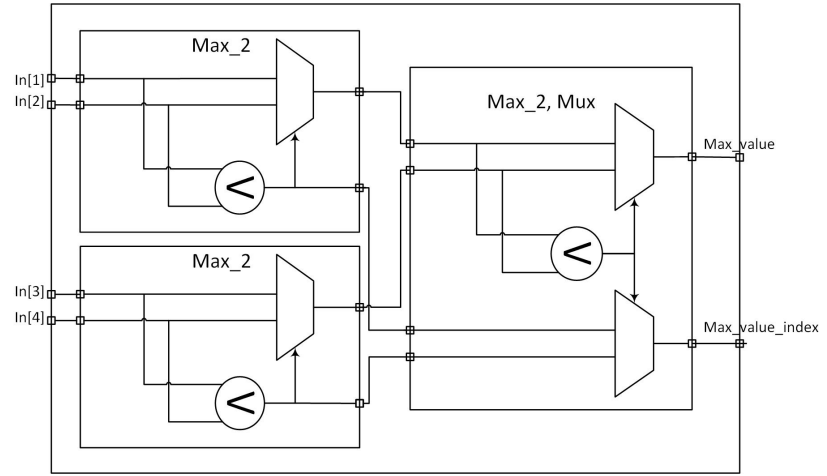Louisiana State University   Alibaba Group
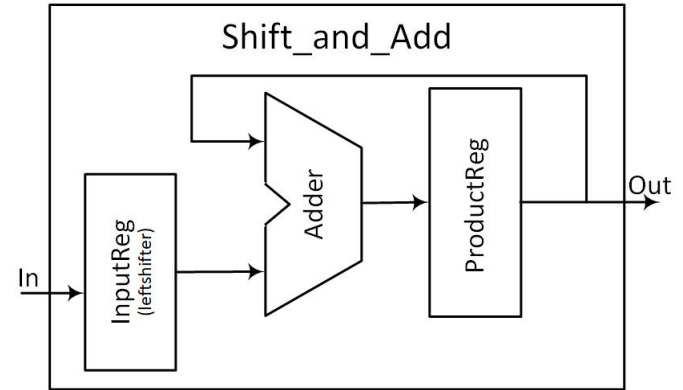
A³ design

A³ architecture

Memristor crossbar

Max pooling unit

Shift & Add unit

A Virtual Image Accelerator for Graph Cuts Inference on FPGA

Tianqi Gao, Rob A Rutenbar

Graph Cuts Inference via Push Relabel

Push(v, w).
Applicability: v is active, $r_f(v, w) > 0$ **and** $d(v) = d(w) + 1$.
Action: Send $\delta = min(e(v), r_f(v, w))$ units of flow from v to w
$f(v, w) \leftarrow f(v, w) + \delta; f(w, v) \leftarrow f(w, v) - \delta;$
$e(v) \leftarrow e(v) - \delta; e(w) \leftarrow e(w) + \delta.$

Relabel(v).
Applicability: v is active **and** $\forall w \in V, r_f(v, w) > 0 \Rightarrow d(v) \le d(w)$.
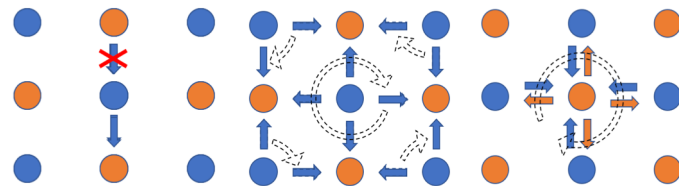Action: $d(v) \leftarrow min\{d(w) + 1 \mid (v, w) \in E_f\}.$
(If this minimum is over an empty set, $d(v) \leftarrow \infty$.)

System Design: Pixel-Parallel Architecture

- Design a physical tile of pixel-processors to execute Push Relabel in parallel
- Design a Virtual-Image system for large images by dividing it into *virtual tiles*
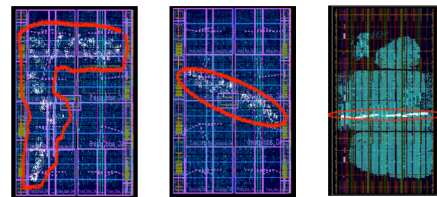
Single Tile

Checkerboard Scheduling

Implement 2304 pixel-processors

Placement
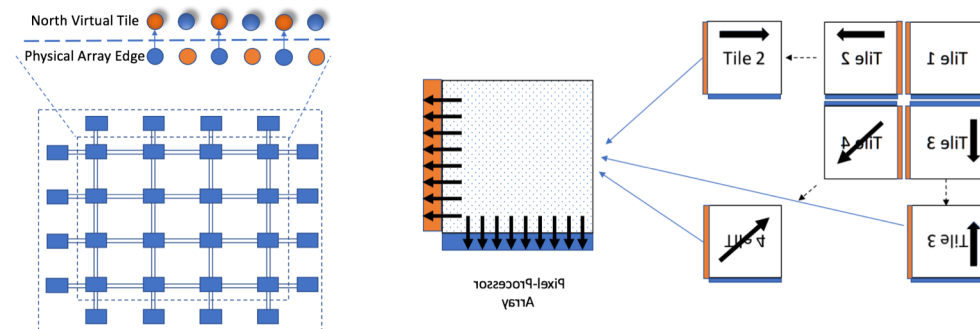- Adjust the weight-height ratio to fit on the FPGA

Routing
- Use a global distribution tree

Virtual-Image Architecture

- Memory Virtualization
- Shadow Edge Processors for Flow between Virtual Tiles
- Properly Mapping Virtual Tiles onto the Physical Tile

Experiments Results

- 11-13x faster than other FPGAs
- 1.38x faster than a GPU implementation

Benchmark images (first row) used in our experiments with their results (second row).

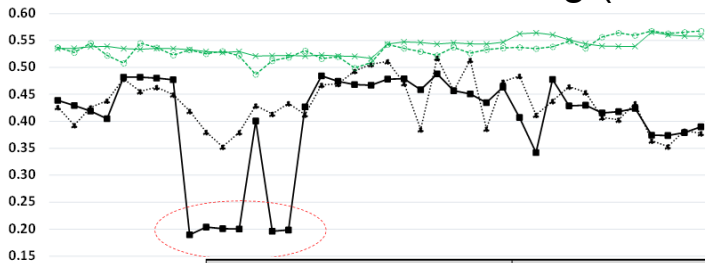# Implications for Hardware Acceleration of Malware Detection

Jordan Pattee and **Byeong Kil Lee**

{ jpattee, blee } @uccs.edu

University of Colorado
Colorado Springs

University of Colorado
Boulder | Colorado Springs | Denver | Anschutz Medical Campus

- ❑ Software-based malware detection
  → **Hardware-based Solution**

- ❑ **PMC**s (performance monitoring counters)
  - ✓ limited number of PMCs in a processor
  - ✓ impact the performance of host processor

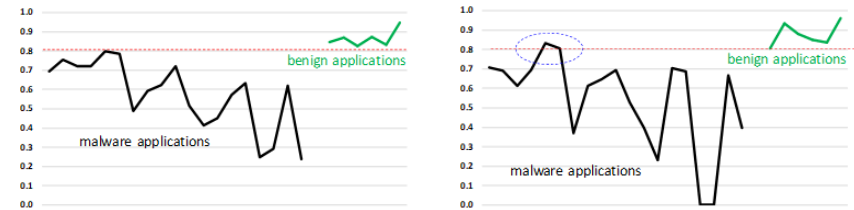- ❑ A metric to differentiate malwares from benign applications (DoD)

$$Degree\ of\ Distribution = \frac{Mean}{(Mean + StdDev)}$$

- ❑ **Feature Selection** for Learning (40 vs. 6)

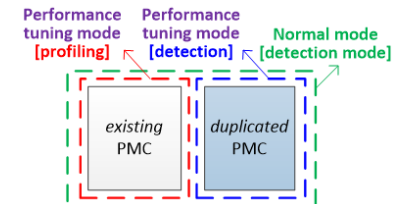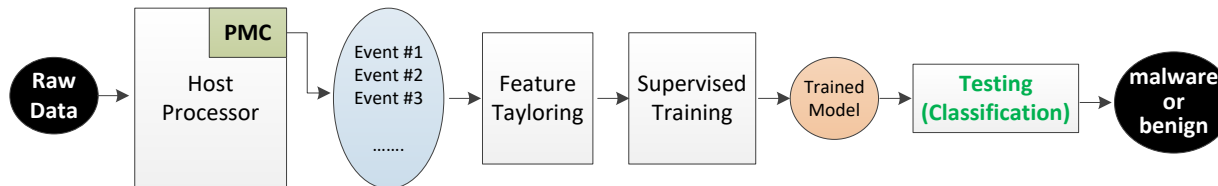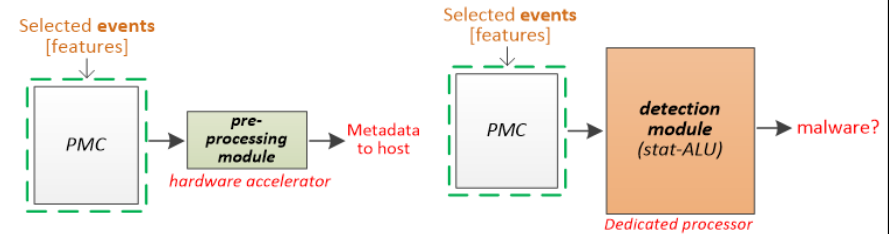| Degree of Distribution (DoD)	Attribute Evaluation
L1D Read Accesses	Cache References
L1D Read Misses	Cache Misses
L1D Write Accesses	Branch Instructions
L1D Write Misses	Branch Misses
L1I Prefetch Accesses	Bus Cycles
L1I Read Accesses	L1D Read Accesses

- ❑ **Malware Detection (40 vs. 6)**

- ❑ **Performance Comparison**

	False Positive	True Negative	F-measure	AUC (ROC)	AUC (PRC)
6-attrib-standard	0.44	0.56	0.93	0.85	0.92
6-attrib-3-fold	0.46	0.54	0.93	0.85	0.94
6-attrib-5-fold	0.42	0.58	0.94	0.86	0.98
6-DoD-standard	0.15	0.85	0.97	0.99	1.00
6-DoD-3-fold	0.26	0.74	0.95	0.96	0.98
6-DoD-5-fold	0.25	0.75	0.96	0.97	1.00

AUC (ROC): Area Under Curve - ROC (Receiver Operating Characteristic) curve
AUC (PRC): Area Under Curve - PRC (associated Precision/ReCall) curve
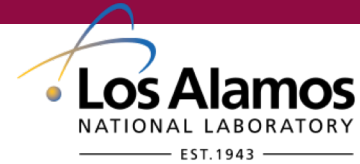
- ❑ **Hardware Acceleration Module**

GPUs Pipeline Latency Analysis

Yehia Arafa, Hameed Badawy, Gopinath Chennupati, Nandakishore Santhi, Stephan Eidenbenz

New Mexico State University, Los Alamos National Laboratory
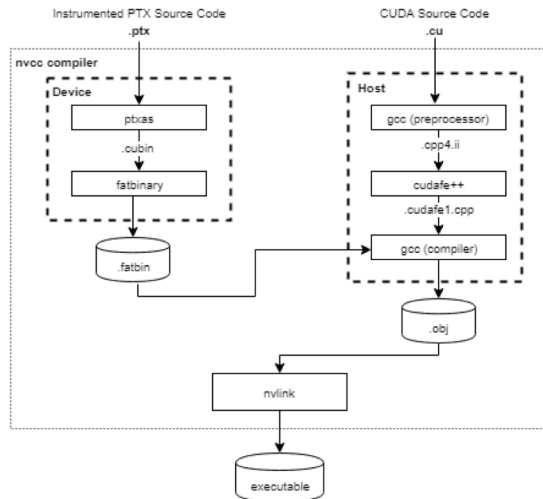
❏ Motivation:

- Graphic Processor Units (GPUs) is now an integral component in any HPC system.
- Over the last decade, Nvidia has introduced seven different generations/architecture. Each has its own microarchitecture and hardware characteristics. However, the percentage of undisclosed characteristics beyond what vendors provides is small.

❏ Goals:

1. Demystify the latency of different instructions executing in the pipeline and different memory units found in various NVIDIA GPUs.
2. Show the effect of high level optimizations found in CUDA (nvcc) compiler on various on the execution of different instructions.

❏ Methodology:

Parallel thread execution (PTX) is used to perform the analysis. To determine each operation's latency, we read the clock register before and after the execution of the instruction.
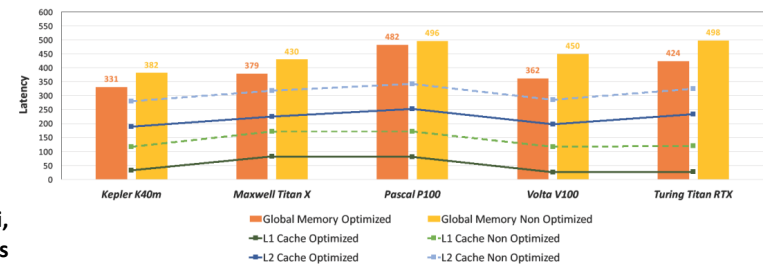

❏ The results are in poster # 93 and in here:

Y. Arafa, A. H. Badawy, G. Chennupati, N. Santhi, and S. Eidenbenz, "Instructions' latencies characterization for nvidia gpgpus," 2019. [Online]. Available: https://arxiv.org/abs/1905.08778

❏ Results:

- We run the evaluation on seven different high-end GPUs from five different generations.
- The results show that the instructions' overhead latencies have mostly decreased from Kepler to Turing.
- We believe that these results should help architects and programmers optimize both the hardware and the software.

Context-Aware Number Generator for Deterministic Bit-stream Computing

Sina Asadi and M. Hassan Najafi

- **Stochastic Computing**
  - Processes data based on bit-streams
  - **Low-cost**, but **low accuracy**, **long latency**, and **high energy consumption**

- **Deterministic Techniques**
  - Introduced recently for completely accurate computation using stochastic logic
  - **Clock Division, Rotation, Relatively Prime,** and **Sobol sequences**
  - **Low-cost** and **high accuracy**, but **long latency**, **high energy consumption**
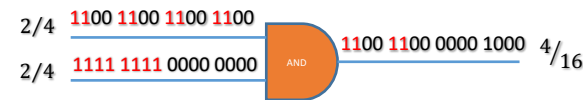
- **Proposed Context-Aware Design**
  - Determines the **minimum** bit-width to precisely represent each input value
  - Includes a proposed small **control unit** and a redesigned architecture
  - **Low-cost, high accuracy**, **faster**, and **lower energy consumption**

- **Evaluation**
  - **Significant improvement in the performance and energy consumption**
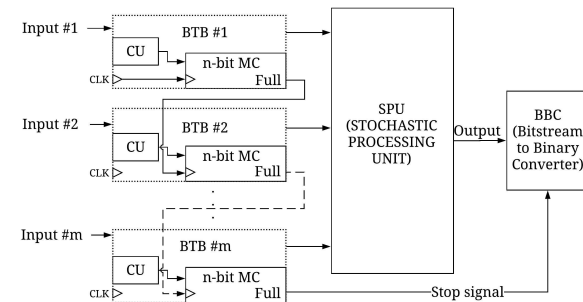 - **Negligible hardware cost overhead**

Multiplication Using Conventional Clock Division Technique

Multiplication Using Proposed Clock Division Technique

Context-Aware Architecture

ASAP 2019

UNIVERSITY of LOUISIANA LAFAYETTE

Smart Rabbit： A Wearable Device As An Intelligent Pacer for Marathon Runners

Yuan Ze University
Wenpei Zheng
Sheng-Yang Chiu
Jui-Chien Hsieh
Chaochang Chiu

Marathon Runner

Sports Therapy & Rehabilitation

Fitness Enthusiast

- **Exercise intensity can be evaluated by ones heart-rate. "Smart Rabbit" helps user to maintain in a specific heart-rate range, aids runners in achieving a certain exercise intensity, preventing sports injuries.**

LSTM RNN NET

AI Heart Rate Predictor

AI心跳預測系統

心跳上限: 108

心跳下限: 92

應用配置　　開始預測　　停止預測

心率正常

横坐標為Time(sec)，纵坐標為心率HeartRate(bpm)

心率低警戒線

— 实际心率 (bpm)　— 预测心率 (bpm)

Polar H7 75D...　DISCONNECT

131 bpm
0.91 km
1.0 m/sec
3.7 mets
00:05:27